Per-Mount Ownership Changes

Christian Brauner christian@brauner.io @brau_ner

Ownership

- uids and gids express ownership
- VFS uses them for permission checking
- persisted to disk for FS_REQUIRES_DEV filesystems

Ownership & struct inode

- i_uid_read()
 - read ownership information from struct inode
 - calls from_kuid() to translate kuids to raw uids
- i_uid_write()
 - write ownership information to struct inode
 - calls make_kuid() to translate raw uids into kuids

ID mappings

- translation of range of ids into another or same range of ids
- notational convention in this talk: u:k:r
 - u == userspace-id / userspace-idmapset
 - k == kernel-id / kernel-idmapset
 - r == range
- associated with struct user_namespace
- init_user_ns has identity idmapping: u0:k0:r4294967295

ID mappings

```
make_kuid(u0:k10000:r10000, u1000)
  What does u1000 map down to?
  id - u + k = n
  u1000 - u0 + k10000 = k11000
- from_kuid(u0:k10000:r10000, k11000)
  What does k11000 map up to?
  id - k + u = n
  k11000 - k10000 + u0 = u1000
```

Ownership: Disk to VFS

file owned on disk by raw uid 1000

```
- fs mounted in init_user_ns
   i_uid_write(u0:k0:r4294967295, u1000) = k1000
```

- fs mounted with idmapping
 i_uid_write(u0:k10000:r10000, u1000) = k11000

```
// Examples
xfs_inode_to_disk(), ext4_do_update_inode(), fill_inode_item() // btrfs
```

Ownership: VFS to Disk

- file owned on disk by raw uid 1000

```
- fs mounted in init_user_ns
i_uid_write(u0:k0:r4294967295, u1000) = k1000
i_uid_read(u0:k0:r4294967295, k1000) = u1000
```

- fs mounted with idmapping

```
i_uid_write(u0:k10000:r10000, u1000) = k1100
i_uid_read(u0:k10000:r10000, k11000) = u1000
```

```
// Examples
xfs_inode_from_disk(), __ext4_iget(), btrfs_read_locked_inode()
```

Creating New Files (Userspace to/from VFS)

Translate between two ID-mappings via the kernel idmapset:

- Map caller's userspace ids down into kernel ids in the caller's idmapping.
 // current_fsuid()
- Verify caller's kernel ids can be mapped up to userspace ids in filesystem's idmapping.

```
// fsuidgid_has_mapping()
```

Crossmapping

Filesystem-wide Idmappings

- alter ownership filesystem-wide
- relevant idmapping is represented in the filesystem's superblock
- determined at mount time

Selected Modern Filesystem Use-Cases

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- Portable Home Directories
 - make login uid and gid random
 - take home directory between computers

Selected Modern Filesystem Use-Cases

- Portable Home Directories
 - make login uid and gid random
 - take home directory between computers
- Containers
 - rootfs
 - data sharing host <> container
 - data sharing container <> container

File ownership should be changeable on a per-mount basis instead of a filesystem wide basis.

Idmapped mounts make it possible to change ownership in a temporary and localized way:

- ownership changes are restricted to a specific mount
- ownership changes are tied to the lifetime of a mount

Idmapping functions were added that translate between idmappings:

make_kuid(mount, uid) = kuid

- i_uid_into_mnt()*
 - translate filesystems kernel ids into kernel ids in the mount's idmapping

```
/* Map filesystem's kernel id up into a userspace id in the filesystem's idmapping. */
from_kuid(filesystem-idmapping, kid) = uid
/* Map filesystem's userspace id down into a kernel id in the mount's idmapping. */
```

- mapped fsuid()
 - translate caller's kernel ids into kernel ids in the filesystem's idmapping by remapping the caller's kernel ids using the mount's idmapping

```
/* Map the caller's kernel id up into a userspace id in the mount's idmapping. */
from_kuid(mount-idmapping, kid) = uid

/* Map the mount's userspace id down into a kernel id in the filesystem's idmapping. */
make_kuid(filesystem-idmapping, uid) = kuid
```

*In our documentation I call it "remapping algorithm" because it undoes an existing idmapping and remaps it according to the mount's idmapping.

Idmapped Mounts: Portable Home Directories

Verify that the caller's kernel ids can be mapped to userspace ids in the filesystem's idmapping

from kuid(u0:k0:r4294967295, k1000) = u1000 // VFS to Disk

So ultimately the file will be created with raw uid 1000 on disk.

vfs_mkdir()

3.

Idmapped Mounts: Portable Home Directories

```
vfs_getattr() + cp_statx()
      caller id:
                      u1001
      caller idmapping: u0:k0:r4294967295
       filesystem idmapping: u0:k0:r4294967295
      mount idmapping: u1000:k1001:r1
      Map the userspace id on disk down into a kernel id in the filesystem's idmapping
      make_kuid(u0:k0:r4294967295, u1000) = k1000 // i_uid_write()
      Translate the kernel id into a kernel id in the mount's idmapping
       i_uid_into_mnt(k1000)
              /* Map the kernel id up into a userspace id in the filesystem's idmapping. */
              from_kuid(u0:k0:r4294967295, k1000) = u1000
              /* Map the userspace id down into a kernel id in the mounts's idmapping. */
              make_kuid(u1000:k1001:r1, u1000) = k1001
  3.
      Map the kernel id up into a userspace id in the caller's idmapping
```

 $from_kuid(u0:k0:r4294967295, k1001) = u1001 // VFS to Userspace$

So ultimately the caller will be reported that the file belongs to raw uid 1001 which is the caller's userspace id in our example.

```
struct mount_attr *attr = &(struct mount_attr){};
int fd_tree = open_tree(-EBADF, source,
                        OPEN_TREE_CLONE | OPEN_TREE_CLOEXEC |
                        AT_EMPTY_PATH | AT_RECURSIVE);
attr->attr_set |= MOUNT_ATTR_IDMAP;
attr->userns_fd = fd_userns;
mount_setattr(fd_tree, "", AT_EMPTY_PATH | AT_RECURSIVE,
              attr, sizeof(struct mount_attr));
```

Filesystem Support & Adoption

- fat, ext4, xfs, btrfs, ksmbd and more to come
- Already widely adopted in userspace with and a variety of patchsets out there LXD, containerd, systemd, ...
- not a container feature!